Performance Evaluation and Analysis of Delay Tolerant Networking

Earl Oliver, Hossein Falaki

David R. Cheriton School of Computer Science University of Waterloo

11th June 2007

イロト 不同下 イヨト イヨト

3

Motivation

Delay Tolerant Networking

Many projects (KioskNet, DieselNet and SenNDT) use **DTN** over *opportunistic* wireless connections.

Problem Definition

- What are the hardware requirements for DTN applications?
- 2 What is the optimal configuration in terms of performance?
- What methodology should be used to study DTN performance?

Motivation

Delay Tolerant Networking

Many projects (KioskNet, DieselNet and SenNDT) use **DTN** over *opportunistic* wireless connections.

Problem Definition

- **0** What are the hardware requirements for DTN applications?
- **2** What is the optimal configuration in terms of performance?
- What methodology should be used to study DTN performance?

Outline



2 Microbenchmarks

3 Hypotheses



<ロト < 部 > < 言 > < 言 > 言 の < で 3/21

Outline









<ロ > < 部 > < 言 > < 言 > 言 の < C 3/21

Outline









<□> <⊡> <⊡> < 글> < 글> < 글> < 글> ⊇ ♡ < ⊙ 3/21

Outline

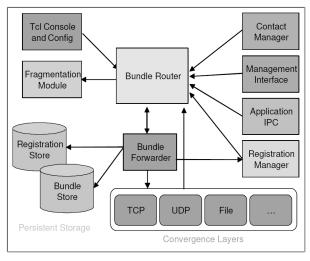








DTN Reference Implementation (DRI)



Methodology

Steps

- We chose common DTN hardware
- Obtermined its capacity with microbenchmarks
- **③** Hypothesized about **DRI performance**
- O Evaluated the hypotheses

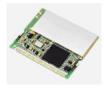
DRI Performance

Throughput between DRI nodes during opportunistic connections.

Test Bed

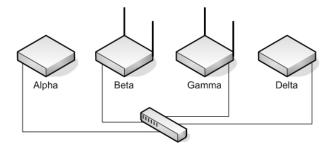
H/S Configuration

- Hardware: Soekris net4801: 266 MHz processor and 256 MB SDRAM
- WiFi card: Atheros 802.11abg wireless cards
- OS: Stable Debian with Linux Kernel 2.6.8-3
- DRI: DTNRG CVS head as of February 22, 2007.





Test Bed Topology



- Wired network as control plane
- Wireless connection between Beta and Gamma

Storage I/O Microbenchmarks

Motivation

Measuring the actual storage throughput limitations

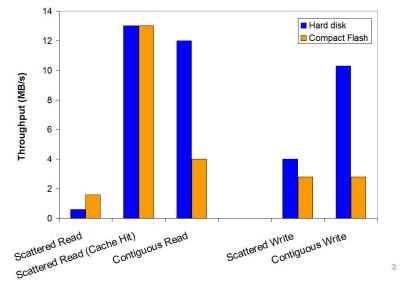
Storage I/O Throughput

- Scattered I/O: reading/writing from/to 1000 to 4000 files
- Contiguous I/O: reading/writing a single large file





I/O Microbenchmarks



Network Microbenchmarks

Motivation

Measuring the actual wireless network throughput limitations

Network I/O Throughput

- Single TCP connection without any disk involvement (3.1 MB/s)
- Reading and writing from/to disk on TCP ends (1.9 MB/s)

Hypotheses: Resource Limitations

• **CPU**: We do not expect to be the primary bottleneck, although it is a scarce resource.

• **Disk**: We expect to be the primary bottleneck of the DRI

• **Memory**: Page faults do not limit the DRI performance





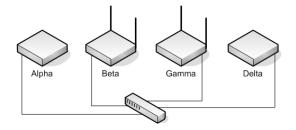
Hypotheses: System and Application Parameters

- Wireless data rate: The DRI performance will peak at a rate of at most 24 Mb/s
- Bundle size: Increasing bundle size will improve performance
- Parallelism: Performance could be improved by parallelism

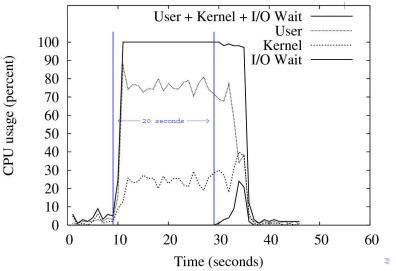
Experiments

Method

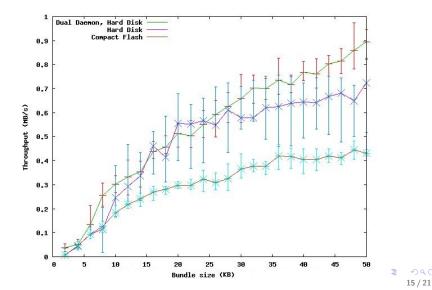
- Enqueue a fixed amount of data at beta
- Simulate a 20 second opportunistic connection
- Measure data remaining in queue after connection is closed



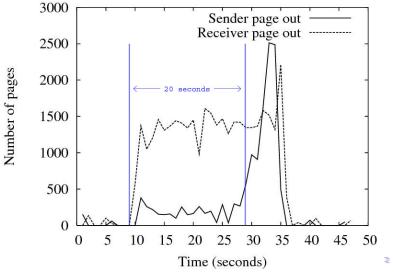
CPU activity during a wireless opportunistic connection



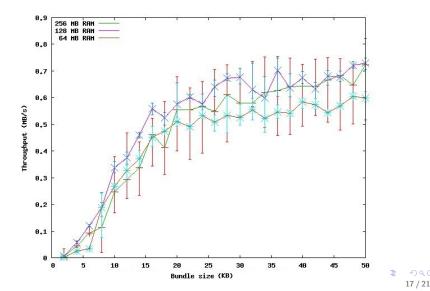
Data throughput vs. bundle size



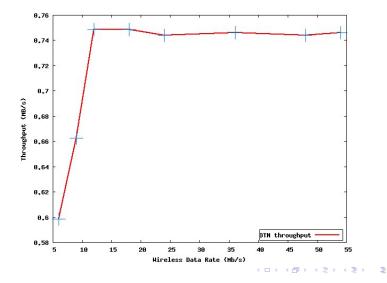
Disk block I/O during a wireless opportunistic connection



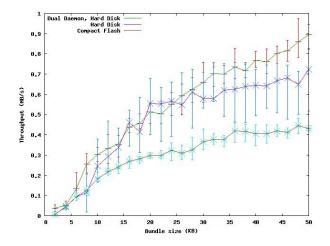
Data throughput vs. bundle size using a portion of memory



Data throughput vs. wireless data rate



Parallelism



• 16% higher throughput (despite being CPU bound)



- 8.28% less time spent on spinlocks
- For example, in "servlib/bundling/BundleList.cc" line 31:

```
BundleList::BundleList(const std::string& name)
        : Logger("BundleList", "/dtn/bundle/list/%s", name.c_str()),
            name_(name), lock_(new oasys::SpinLock()), notifier_(NULL)
{
}
```

Summary

Conclusion

- The primary bottleneck to the DRI performance on common hardware is **CPU**
- Bundle size highly affects DRI performance
- Methodology for evaluating the performance of other mobile systems using opportunistic connections

Recommendations

- **Application developers**: use the largest possible bundle size for the DRI in-memory API
- DRI developers: restructure DRI to increase parallelism or remove spinlocks
- DRI users: invest more in the CPU

Summary

Conclusion

- The primary bottleneck to the DRI performance on common hardware is **CPU**
- Bundle size highly affects DRI performance
- Methodology for evaluating the performance of other mobile systems using opportunistic connections

Recommendations

- Application developers: use the largest possible bundle size for the DRI in-memory API
- DRI developers: restructure DRI to increase parallelism or remove spinlocks
- DRI users: invest more in the CPU